

159201

Week 1

Summer 2014



Data Types

- Basic Data Types
- Integers, real, characters, boolean ...
- C++
 - int
 - float
 - char
 - **bool** (C has no boolean types, programmers use #define etc)

Data Types

- Basic Data Types **grouped together**
- *Structured Data Types:*
- *Arrays, strings, records*
- *In C++ we can use **struct***

Data Types

```
struct BookRecord {  
    char title[40];  
    float callnumber;  
};  
BookRecord book;  
book.callnumber = 5.265;
```

Abstract Data Types (ADT)

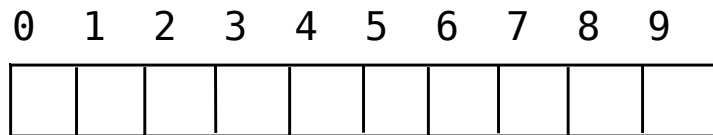
- Specification separate from implementation
- Example:
 - A book record consists of:
 - Title (max 40 characters)
 - Call number (real)

Abstract Data Types (ADT)

- Advantages of ADTs
 - Reduce details – allow focus to be on the “main picture”
 - Different implementations can be used – e.g., array or linked-list
 - Underlying implementation can be changed or upgraded
 - It is convenient to implement an ADT as a **class**

Revision of Arrays

- Remember arrays in C or C++? Example:
 - `int x[10];` // ten elements `x[0]`, `x[1]` ... `x[9]`



These are the **index** numbers



Revision of Arrays

Advantages of Arrays

Simple, Fast, Random access

Disadvantages of Arrays

Every element if of the same data type

Fixed size – too small or too big at runtime

Difficult to insert or delete without leaving spaces



2D arrays

Example:

```
int matrix[4][4];
```

At some point, `matrix[2][2]=64;`

	0	1	2	3
0				
1				
2			64	
3				



2D arrays

- 2D Arrays in C/C++ are stored as a 1D array
 - **Row-major** order
 - Known in math as a **matrix**
 - **Sparse matrix** has few numbers and lots of elements with value = 0



Row-major X column-major

Row-major order? How do we know?

```
#include <stdio.h>
```

```
int a; int b;
```

```
int matrix[4][4];
```

```
main(){
```

```
    for(a=0;a<4;a++){
```

```
        for(b=0;b<4;b++){
```

```
            printf("%ld ",&matrix[a][b]); //pointers
```

```
        }
```

```
        printf("\n");
```

```
    }
```

```
}
```



Row-major X column-major

Output:

6293920 6293924 6293928 6293932
6293936 6293940 6293944 6293948
6293952 6293956 6293960 6293964
6293968 6293972 6293976 6293980

- the output may not be the same for different machines, even for different runs.
- However, it follows a pattern: a space of **4 (bytes)** between elements within the same row.
- The first element in the second column is +4 bytes from the last element in the first row
→ row-major confirmed



Reference and pointers

Remember:

- * a pointer (declare a pointer to any type)
- new allocates memory (equivalent to C `malloc()`)
- & the address of a variable.
- > the element of a pointer (that points to a structure)



Examples with *

```
#include <stdio.h>

main(){
    int a=10;
    int *b;
    b=&a;    //the address is the same
    printf("a=%d and b=%d \n",a,*b);
}
```

Result: a=10 and b=10



Examples with funct(type *&)

```
1  #include <stdio.h>
2  int b;
3  void function1(int *a) { a=&b; }
4  void function2(int *&a) { a=&b; }
5  main(){
6      .   int *a;
7      .   int x=10;
8      .   a=&x;
9      .   printf("a=%d ",*a);
10     .   b=20;
11     .   function1(a);
12     .   printf("a=%d ",*a);
13     .   b=30;
14     .   function2(a);
15     .   printf("a=%d \n",*a);
16 }
```

Result: ? ? ?



Examples with funct(type *&)

```
1  #include <stdio.h>
2  int b;
3  void function1(int *a) { a=&b; }
4  void function2(int *&a) { a=&b; }
5  □ main(){
6      .   int *a;
7      .   int x=10;
8      .   a=&x;
9      .   printf("a=%d ", *a);
10     .   b=20;
11     .   function1(a);
12     .   printf("a=%d ", *a);
13     .   b=30;
14     .   function2(a);
15     .   printf("a=%d \n", *a);
16 }
```

Result is: 10 10 30

NOT: 10 20 30



State of memory at line 7

int b



0x44

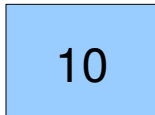
main()

int *a



0x35

int x



0x78

function1(int *A)

int *A



0x97

function2(int *&A)

int *&A



0x89

State of memory at line 9

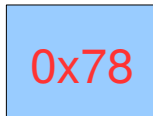
int b



0x44

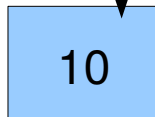
main()

int *a



0x35

int x



0x78

function1(int *A)

int *A



0x97

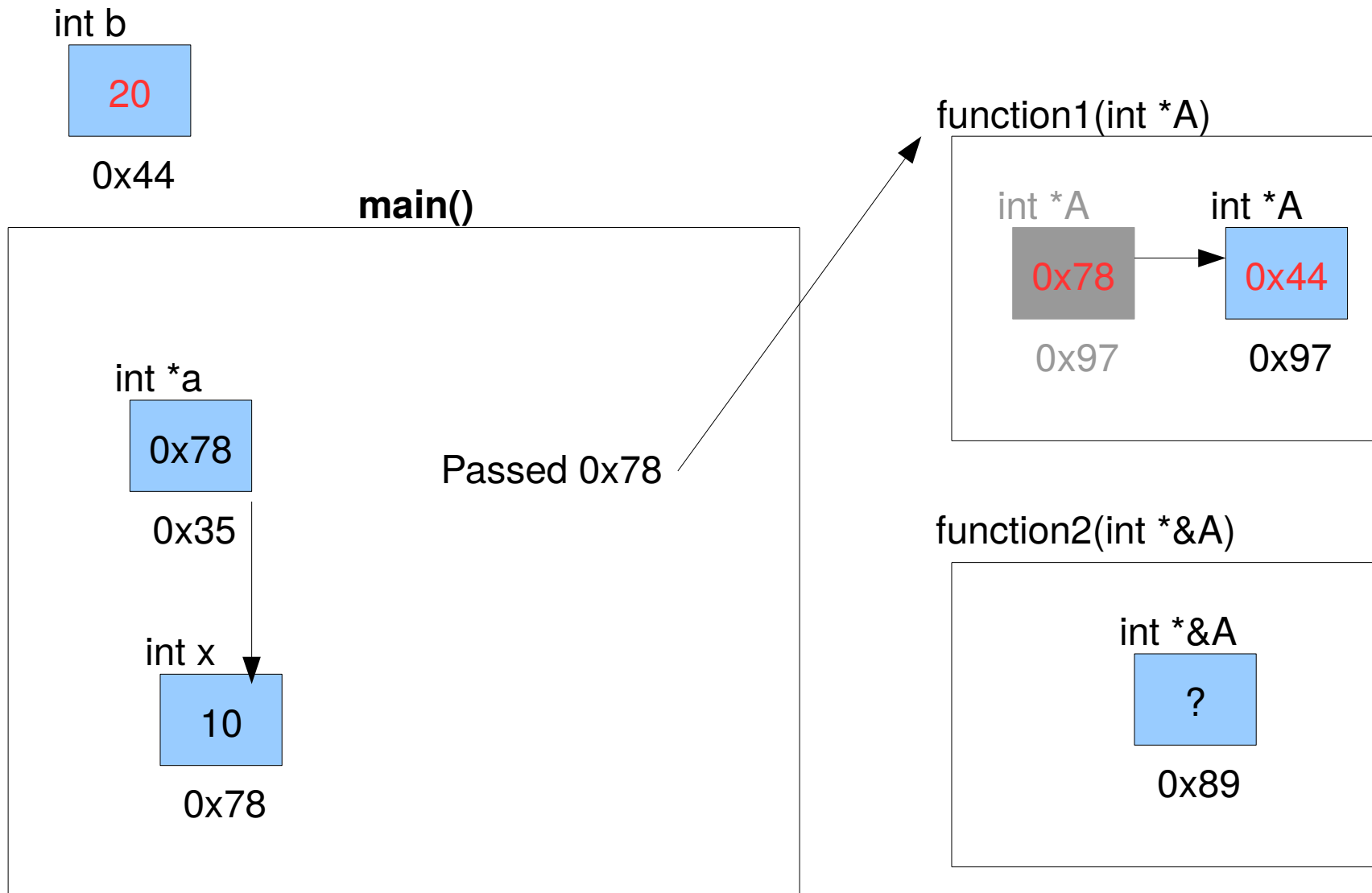
function2(int *&A)

int *&A

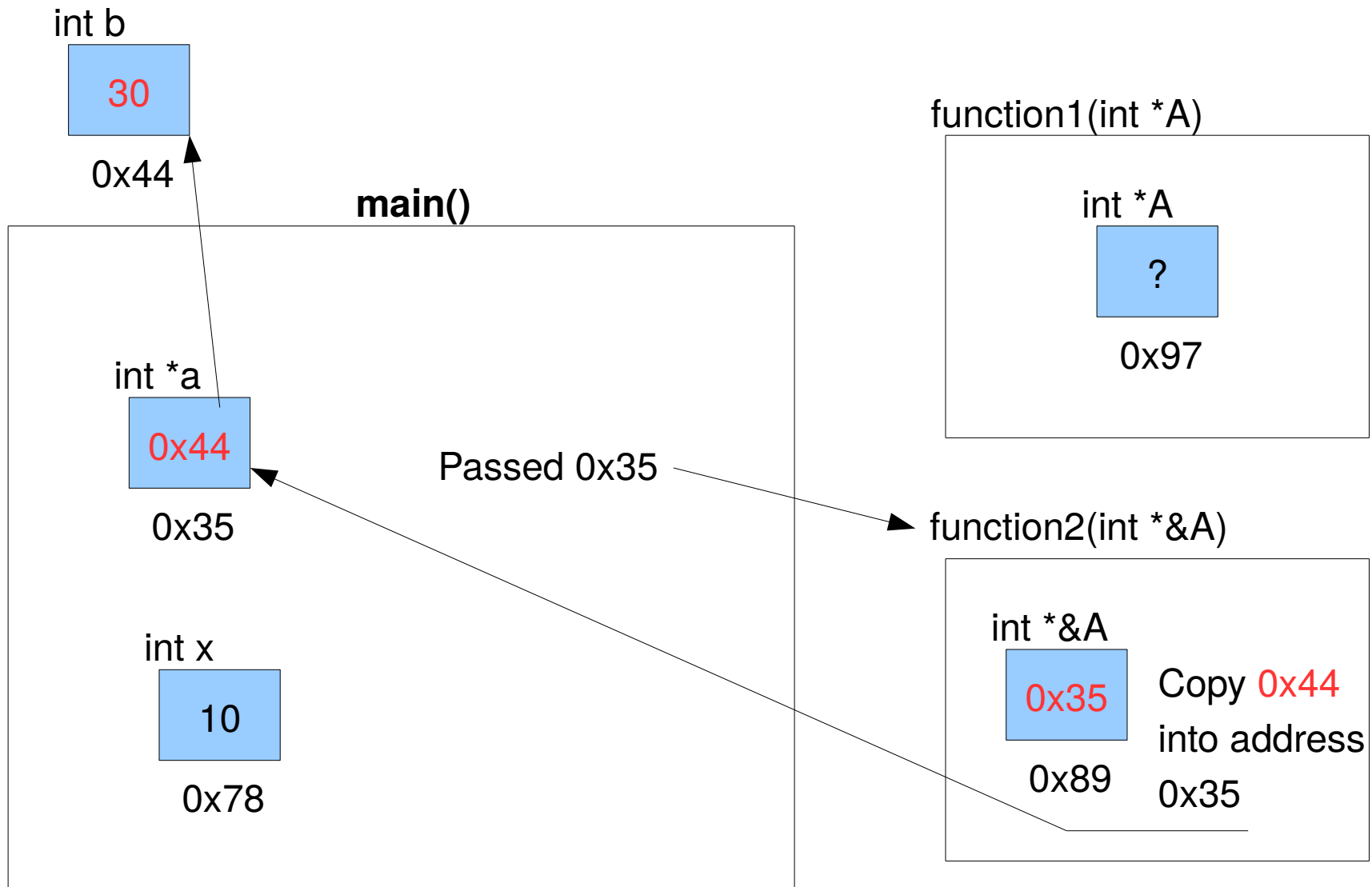


0x89

State of memory at line 12



State of memory at line 15



Examples with `funct(type *&)`

The trick is to pass a pointer to a pointer...

This can be done passing `*&` (typical for C++) or `**` (in C).

Run the program `code1_alternative.cpp` and play with the different variables. Try to follow what is happening to the addresses within the pointers.



malloc() and free(), New, delete

In C, memory allocation/deallocation:

Malloc() and free()

```
#include <stdio.h>
#include <stdlib.h>

main(){
    int a[10];//static, 10 places
    int *b;//pointer only, no allocation yet
    b=(int*) malloc(10*sizeof(int));
    a[5]=10;
    b[5]=10;
    printf("result: a=%d and b=%d\n",a[5],b[5]);
    free(b);
}
```

NOTE: using unallocated pointers or freeing twice leads to disaster... (segmentation fault)

malloc() and free(), New, delete

In C++, memory allocation/deallocation:

New and **delete**

```
1  #include <stdio.h>
2
3  main(){
4      .    int a[10];//static, 10 places
5      .    int *b;//pointer only, no allocation yet
6      .    b = new int[10];//allocate
7      .    a[5]=10;
8      .    b[5]=10;
9      .    printf("result: a=%d and b=%d\n",a[5],b[5]);
10     .    delete[] b;//deallocate (use object's destructor)
11 }
```

NOTE: new and delete have specific roles in OO (constructors and destructors), more in 159234

What -> means?

Remember that “.” is used to refer to elements of structures, e.g.

```
book.callnumber
```

However, when “book” is a pointer we have to refer to it using “->”, e.g.

```
BookRecord book;
```

```
BookRecord *bookpointer;
```

```
...
```

```
Main() {...
```

```
    book.callnumber=10;
```

```
    bookpointer->callnumber=10;
```

```
}
```

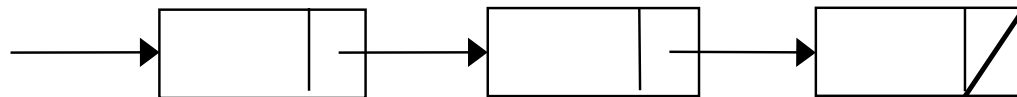


L02



Linked-lists

Linked-lists are sequences of connected nodes.
Linked-lists are empty at the start.
Nodes are added dynamically (at runtime).
Nodes contain pointers to other nodes.
The address of the list is the pointer to the first node.
Linked-lists can be used as an alternative to arrays.



Linked-lists

```
struct Node { //declaration
```

```
    int accnumber;
```

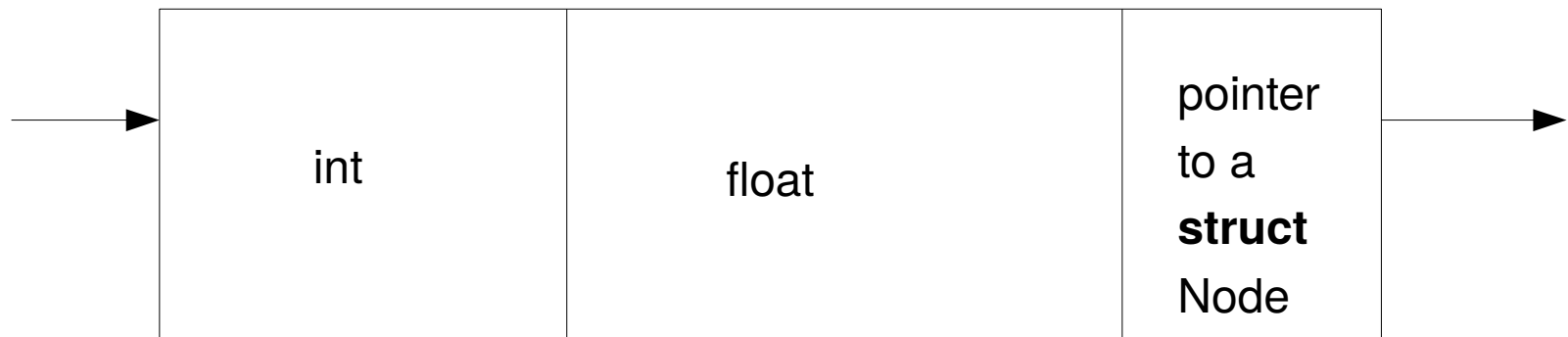
```
    float balance;
```

```
    struct Node *next;
```

```
};
```

```
typedef struct Node Node;
```

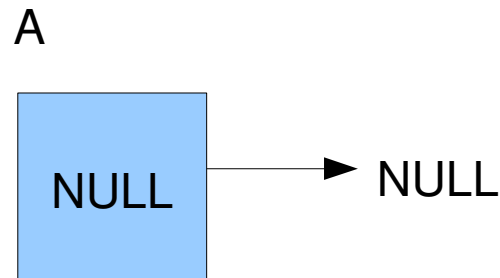
```
//this should reserve memory space for this struct...
```



Linked-lists

Until one declares a Node and specifically allocates memory to it, no memory is allocated:

```
Node *A; //declare one pointer to a linked-list called 'A'  
A = NULL;
```



REMEMBER: there is no place for an int or a float yet...
There is only a pointer to a Node, no allocated memory.

Linked-lists

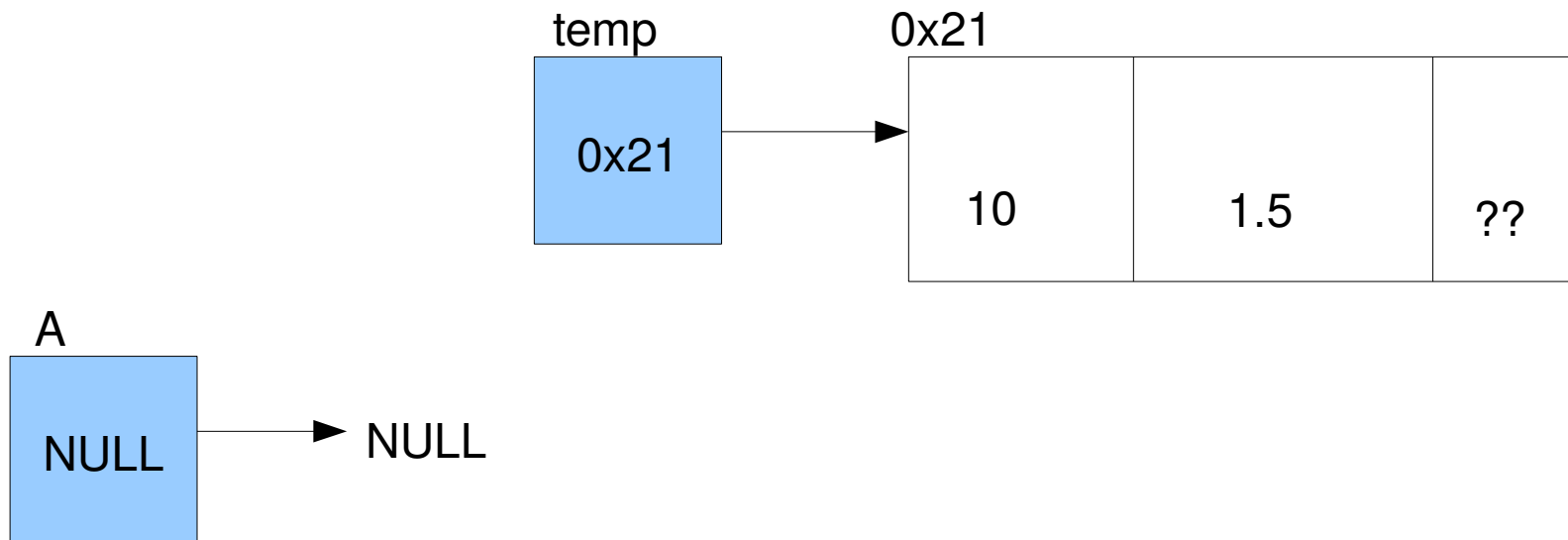
Lets add, manually, a new node on list **A**:

```
Node *temp; //declare a temporary pointer to a Node
```

```
temp = new Node; //allocate space
```

```
temp->accnumber=10; //load the values
```

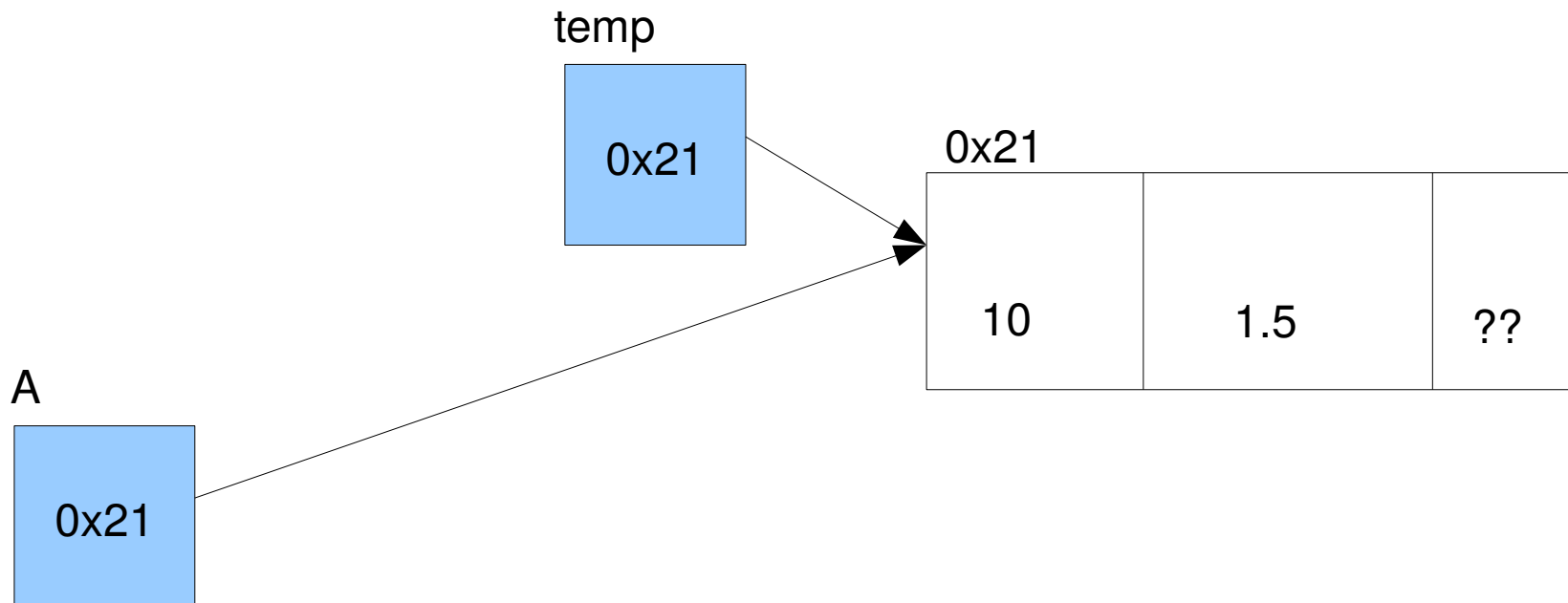
```
temp->balance=1.5;
```



Linked-lists

The new element should be pointed by A. We can copy the content of temp to A:

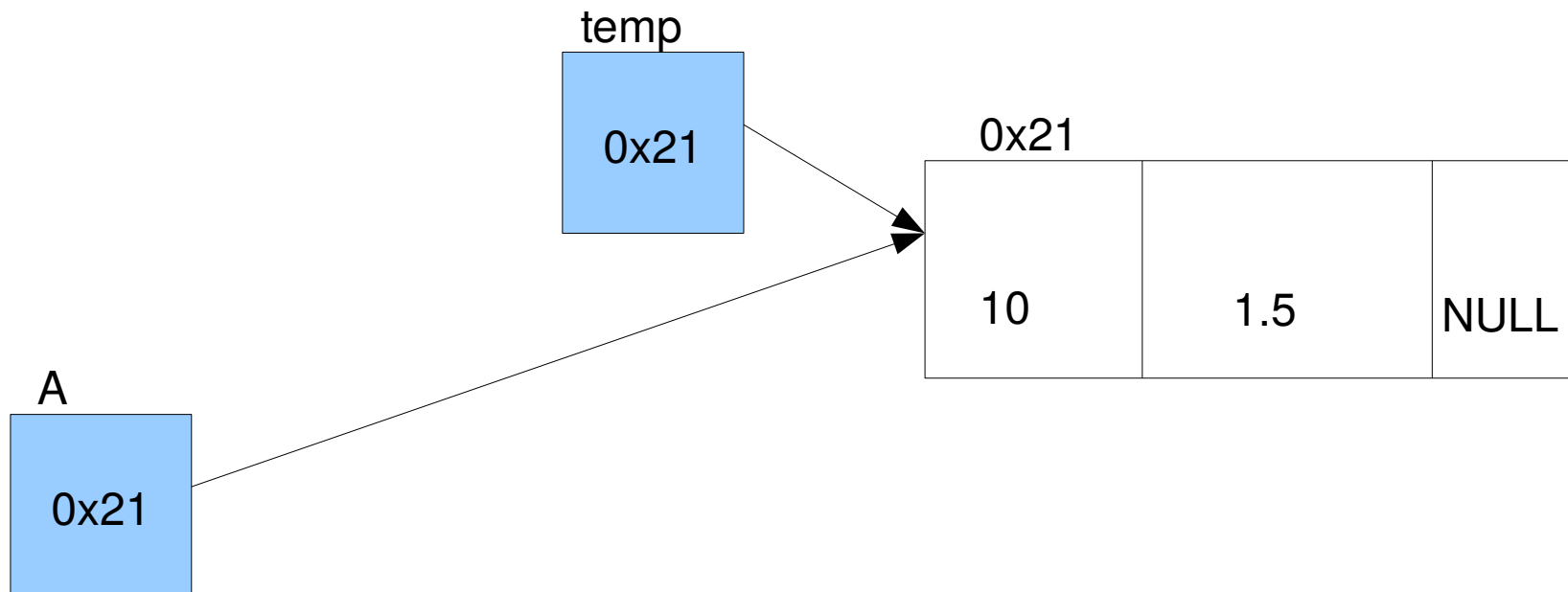
A = temp;



Linked-lists

A has now one element. But the new element points to a **random** place in memory. Lets point it to NULL:

```
temp->next=NULL; //(or A->next=NULL)
```



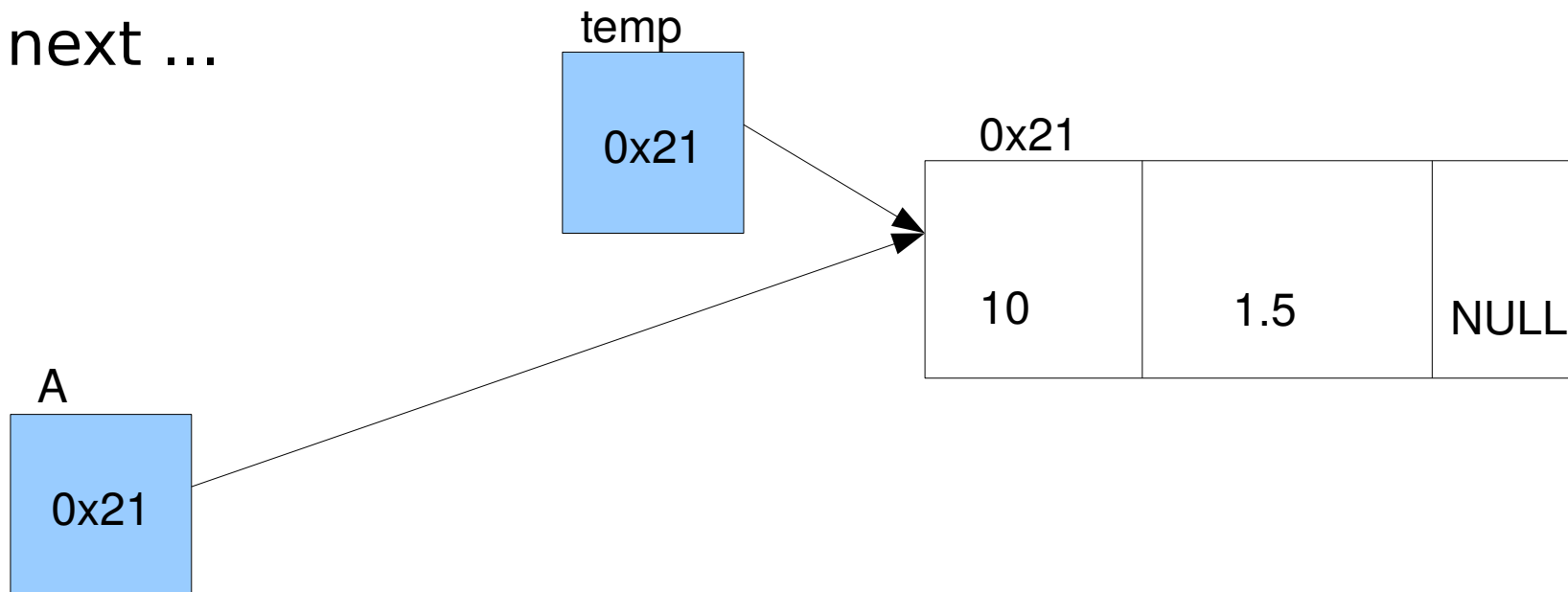
Linked-lists

Now, if we want to refer to the first element of A:

A->accnumber

A->balance

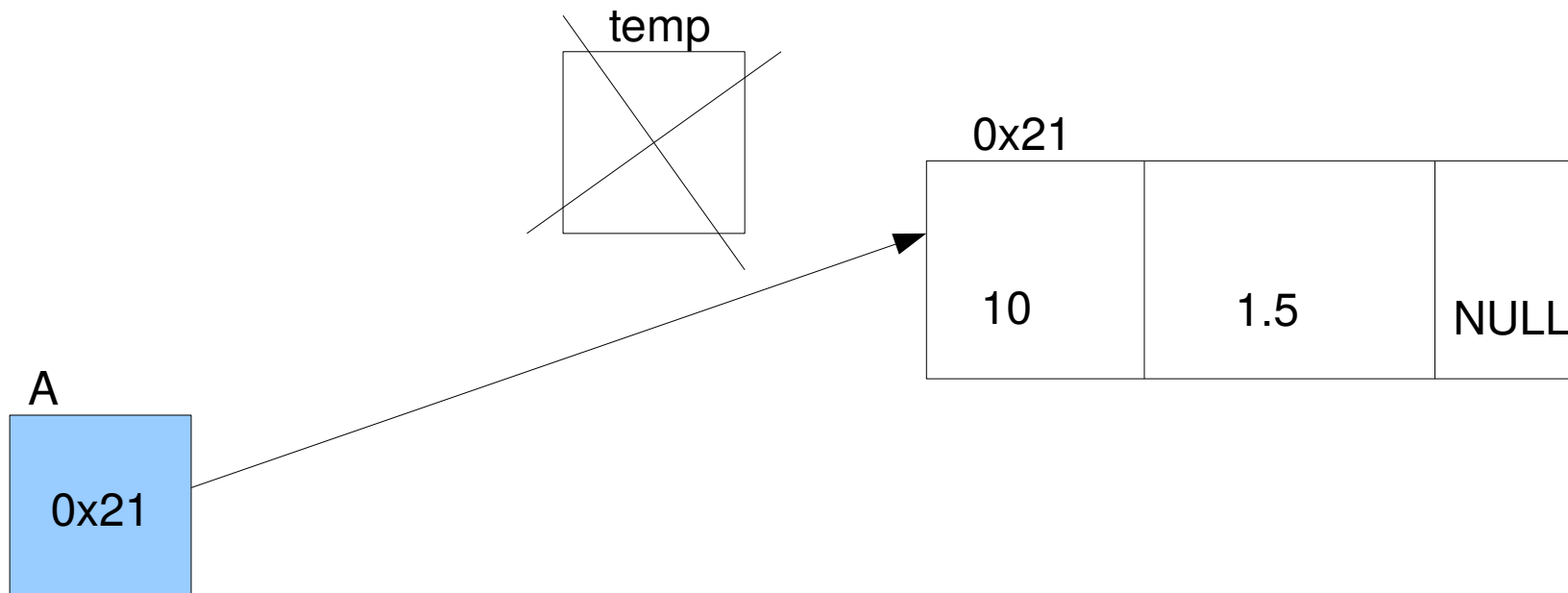
A->next ...



Linked-lists

We could now eliminate `temp` (we will see how to do this properly inside a function)

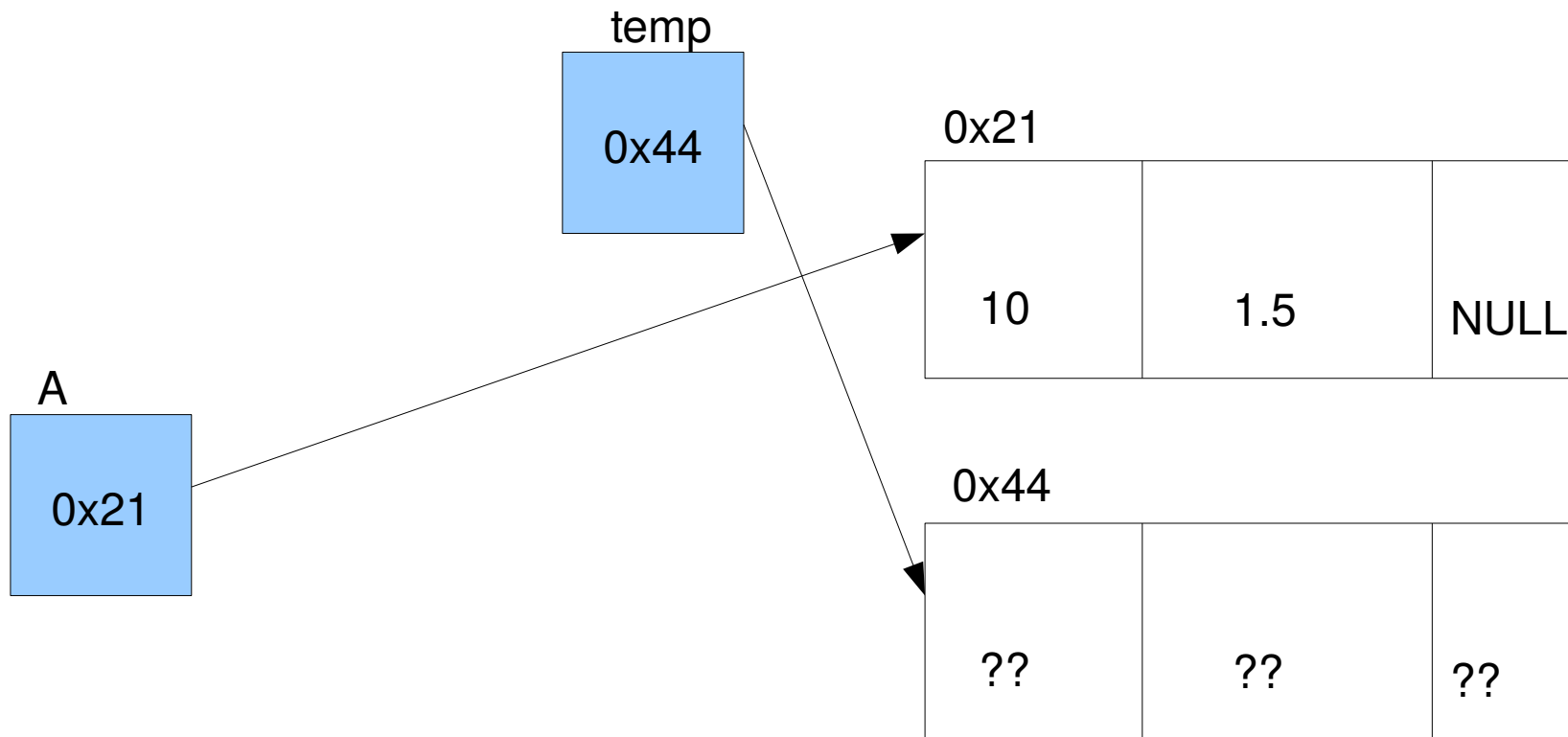
But why do we need `temp` in the first place? Lets use `temp` to create a second element instead of deleting it...



Linked-lists

Suppose you want a second element linked to list A.

```
temp = new Node;
```



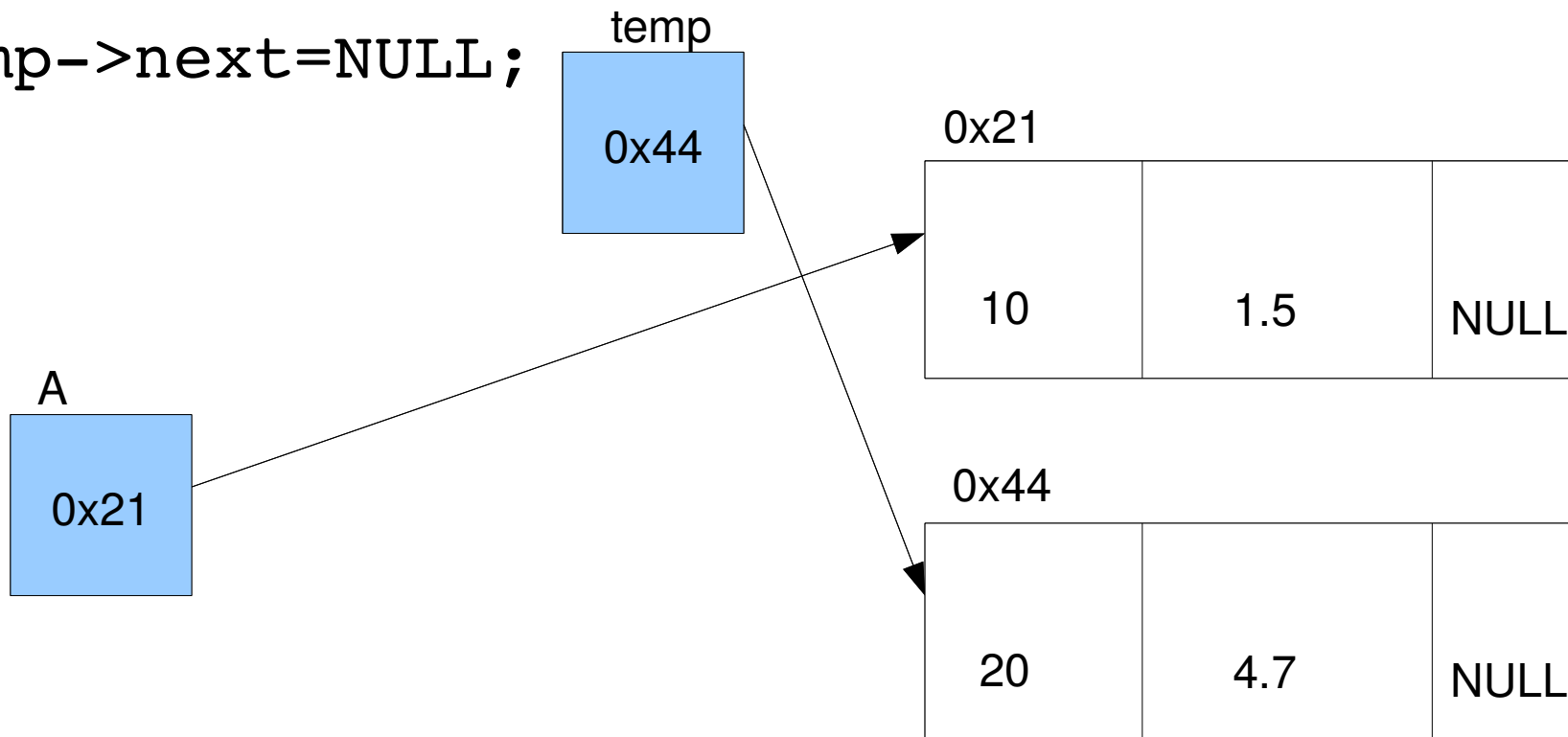
Linked-lists

Load values to it:

```
temp->accnumber=20; //load the values
```

```
temp->balance=4.7;
```

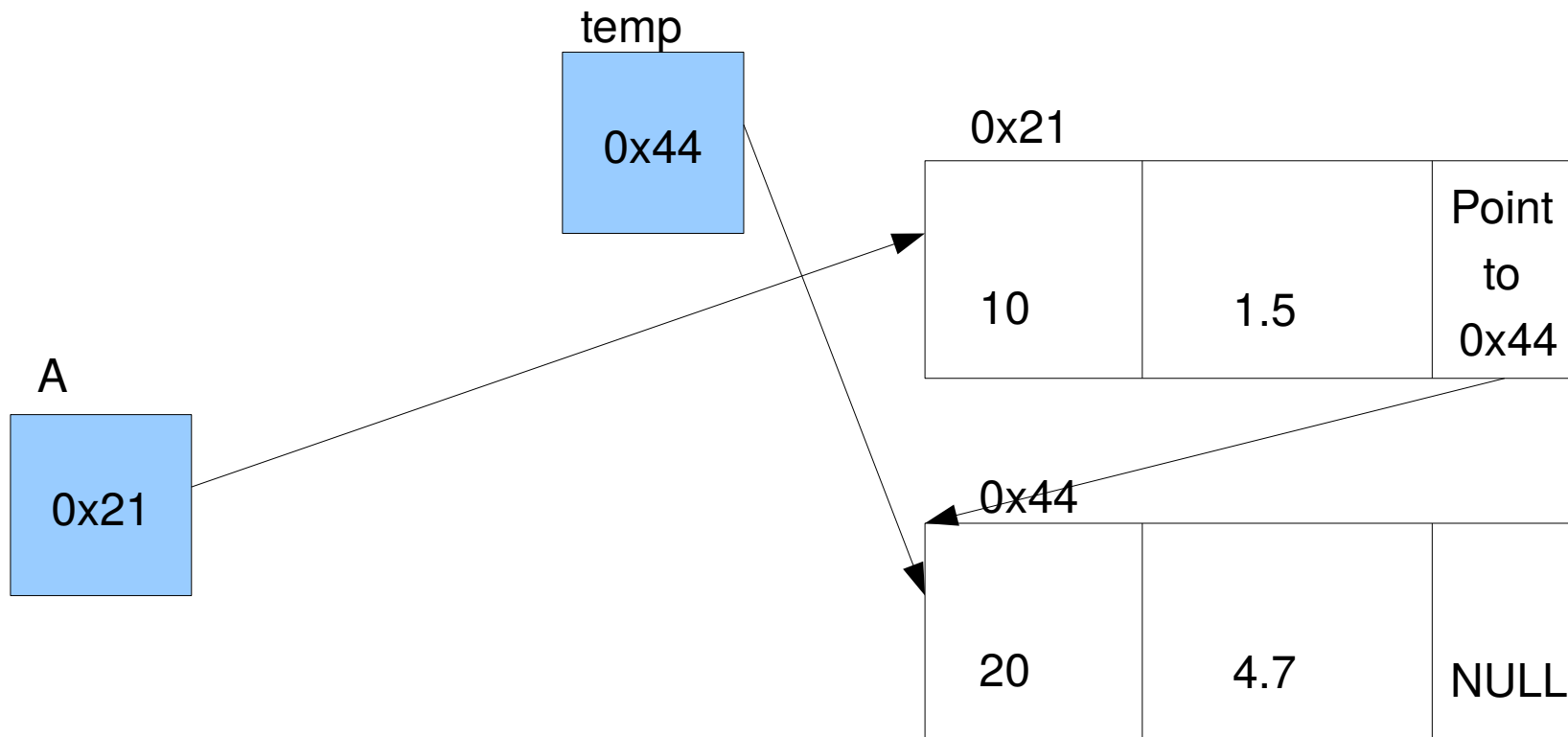
```
temp->next=NULL;
```



Linked-lists

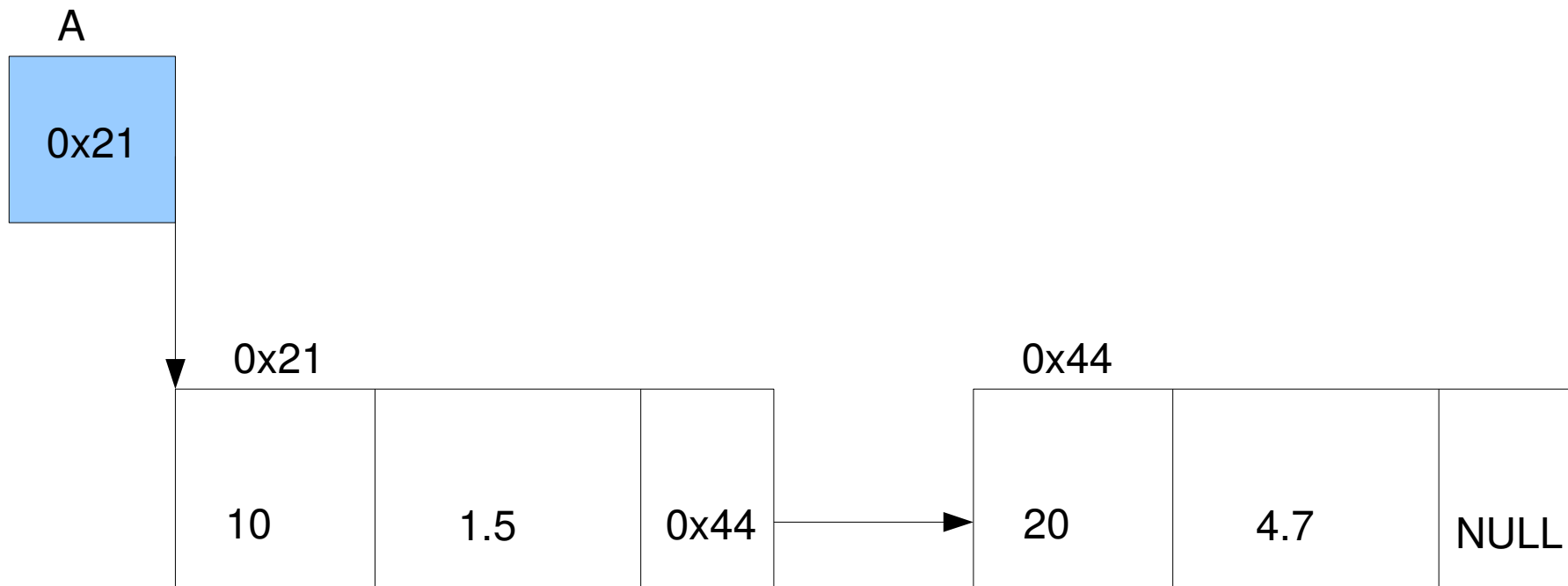
Then, link the second element to the first:

`A->next=temp;`



Linked-lists

Rearranging the figure, this is the state of the linked-list at this point...



Discussion: what happens if we create more elements?

Linked-lists compared to arrays

Linked-lists

Grows during runtime

Dynamic memory allocation

Easy to insert/delete in the middle

Sequential access is fast

slow

Complicated (needs extra functions to work)

Arrays

Fixed size (compilation time)

Static memory allocation

Inserting elements leave empty spaces in memory

Random access (index)

fast

Simple



Sample Linked-list in C++

```
1  #include <stdio.h>
2  struct Node { //declaration
3      int accnumber;
4      float balance;
5      Node *next;
6  };
7  Node *A, *B; //declaration
8
9  void AddNode(Node * & listpointer, int a, float b);
10
11 int main() {
12     A = NULL; // ALL linked-lists start empty
13     B = NULL;
14     AddNode(A, 123, 99.87);
15     AddNode(B, 789, 52.64);
16 }
```

AddNode()

```
19 void AddNode(Node * & listpointer, int a, float b) {  
20     // add a new node to the FRONT of the list  
21     Node *temp;  
22     temp = new Node;  
23     temp->accnumber = a;  
24     temp->balance = b;  
25     temp->next = listpointer;  
26     listpointer = temp;  
27 }  
28
```

But wait a minute...

This will not produce the same linked-list as before, as it will be **inverted** if we add in the same order!

Can you think of two ways of adding nodes, at the HEAD or at the TAIL of the linked list?

Reference and pointers

Subtle syntax in C/C++ can cause errors

A function can get parameters using pointers and/or references:

```
void function1( Node * listpointer...
```

In this case, the pointer to listpointer is passed as reference (a copy of the address is made). Changing listpointer does not alter A or B

```
void function2( Node * &listpointer...
```

In this case, the pointer is itself passed to the function, so changing listpointer changes A or B...



Challenges:

1) Modify the AddNode() function (add to HEAD) to add to the TAIL of the linked-list.

2) Modify the AddNode() function to add an element AFTER a certain element (by value or position of the element).

You will need to find the last element of the linked-list by modifying the Search() function.

The answers are on Stream, study these solutions carefully and understand exactly how to control Nodes: *add, delete and search for any Node.*



L03



Linked-lists Search and Remove

We know how to add nodes to our lists, but just adding to the HEAD. How do we find the TAIL of a linked list?

We also need some extra functions to deal with elements, such as Search. Also, a function to delete or remove nodes that we no longer need.

We need to deal with pointers appropriately to achieve that, it is easy to make a subtle mistake and crash...



Linked-lists Search

Search function:

```
void Search(Node *listpointer, int x) {  
    // search for the node with account number equal to x  
    Node *current;  
    current = listpointer;  
    while (true) {  
        if (current == NULL) { break; }  
        if (current->accnumber == x) {  
            printf("Balance of %i is %1.2f\n", x, current->balance);  
            return;  
        }  
        current = current->next;  
    }  
    printf("Account %i is not in the list.\n", x);  
}
```

Linked-lists Search

Search function:

```
int main() {  
    A = NULL; // ALL linked-lists start empty  
    AddNode(A, 1, 9.87);  
    AddNode(A, 2, 8.87);  
    AddNode(A, 3, 7.87);  
    Search(A, 123);  
    Search(A, 1);  
    Search(A, 2);  
    Search(A, 3);  
}
```

Account 123 is not in the list.

Balance of 1 is 9.87

Balance of 2 is 8.87

Balance of 3 is 7.87

Search step-by-step

- Create a pointer `current`, of same type as node
- `current` initially points to the list, which is the first element of the linked-list
- At any point, if `current` is NULL → reached the **end** of the list (last element)
- We keep checking for `accnumber` and update `current=current->next;`
- Note that we go through the entire list, and we either find the `accnumber` we look for or reach the end of the list



Linked-lists Remove nodes

RemoveNode function:

```
void RemoveNode(Node * & listpointer, int x) {  
    // remove the node containing account number x  
    Node *current, *prev;  
    current = listpointer;  
    prev = NULL;  
    while (current != NULL) {  
        if (current->accnumber == x) { break; }  
        prev = current;  
        current = current->next;  
    }  
    if (prev == NULL) {  
        listpointer = listpointer->next;  
    } else {  
        prev->next = current->next;  
    }  
    delete current;  
}
```

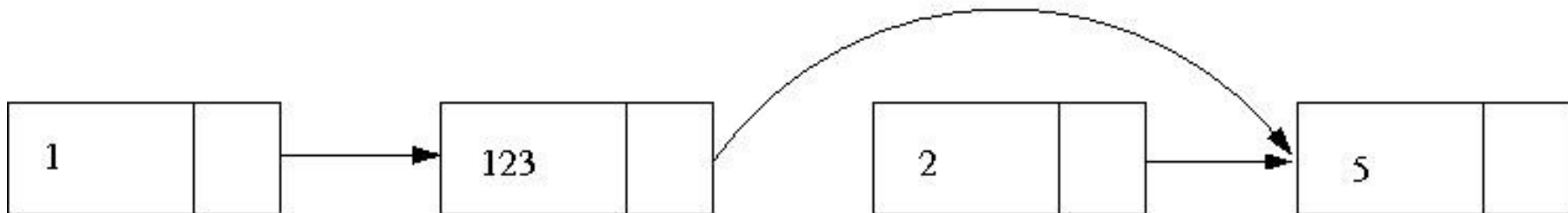

RemoveNode step-by-step

- Two pointers, "current" and "prev"
- current initially points to the list
- prev initially points to nothing (NULL)
- While current is not NULL, search the list until find X. Keep swapping `prev = current`
- If X is found, change prev pointer to jump one element
- Now we can delete the element by deallocating current



RemoveNode in pictures

Suppose we want to remove element
`accnumber==2`:



RemoveNode example

Usage

```
int main() {  
    A = NULL; // ALL linked-lists start empty  
    AddNode(A, 1, 9.87);  
    AddNode(A, 123, 8.87);  
    AddNode(A, 2, 7.87);  
    AddNode(A, 5, 7.87);  
    Search(A, 2);  
    RemoveNode(A, 2);  
    Search(A, 2);  
}
```

Balance of 2 is 7.87

Account 2 is not in the list.

Question!

What happens if:

```
int main() {  
    A = NULL; // ALL linked-lists start empty  
    AddNode(A, 1, 9.87);  
    AddNode(A, 123, 8.87);  
    AddNode(A, 2, 7.87);  
    AddNode(A, 5, 7.87);  
    Search(A, 2);  
    RemoveNode(A, 2);  
    Search(A, 2);  
    RemoveNode(A, 2);  
}
```

Question!

What happens if:

```
int main() {  
    A = NULL; // ALL linked-lists start empty  
    AddNode(A, 1, 9.87);  
    AddNode(A, 123, 8.87);  
    AddNode(A, 2, 7.87);  
    AddNode(A, 5, 7.87);  
    Search(A, 2);  
    RemoveNode(A, 2);  
    Search(A, 2);  
    RemoveNode(A, 2);  
}
```

Balance of 2 is 7.87

Account 2 is not in the list.

Segmentation fault!!!!

Extra pointers

Pointers can be added to point to

rear

middle

one third etc...

Or a combination of the above

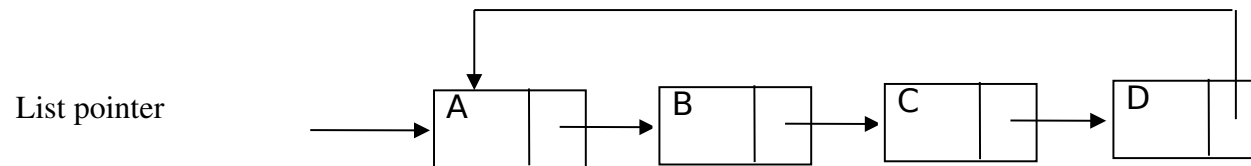
Extra operations on the AddNode() and RemoveNode()

New search functions can be devised. What is the advantage?



Other types of linked-lists

Circular lists



Doubly-linked-lists



L04



Print all elements of a LL

Simple approach: scan LL until the end

```
void PrintLL(Node *listpointer) {  
    // print all elements  
    Node *current;  
    current = listpointer;  
    int element=1;  
    while (true) {  
        if (current == NULL) { break; }  
        printf("Element %d: Balance of acc %i is %1.2f\n",  
            element, current->accnumber, current->balance);  
        current = current->next;  
        element++;  
    }  
    printf("End of the list.\n");  
}
```



More operations with Linked-lists

Extra Operations:

Concatenate → join two separate lists

Reverse → invert the order of the elements

Split → separate the list in two

Insert a new node **after** a certain element

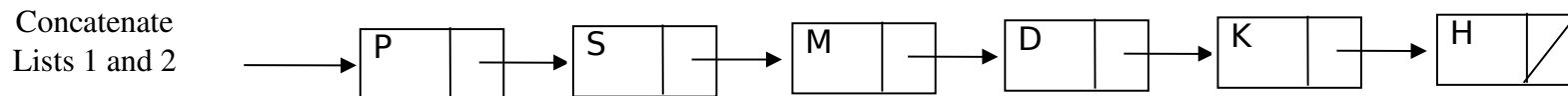
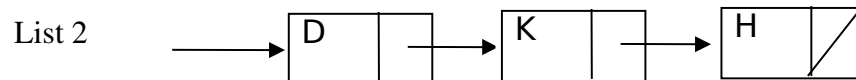
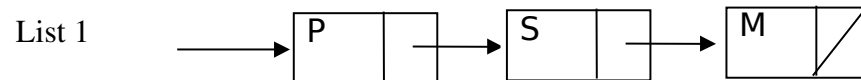
Delete by element **order** (say, the 5th element)
rather than by a known key



Concatenate

Concatenate (join two separate lists)

The final pointer of list 1 should now point to list 2 first element



Concatenate example

- Scan listpointer1 to find the last element
- Join

```
void Concatenate(Node * &listpointer1, Node * listpointer2) {  
    //find the last element of listpointer1, then join with listpointer2  
    Node *current, *prev;  
    current = listpointer1;  
    prev = NULL;  
    while (current != NULL) {  
        prev = current;  
        current = current->next;  
    }  
    if (prev == NULL) {  
        //in this case listpointer1 is empty  
        printf("list1 was empty, join anyway\n");  
        listpointer1 = listpointer2;  
    } else {  
        //join lists  
        printf("join\n");  
        prev->next=listpointer2;  
    }  
}
```

Concatenate example

Main

```
int main() {  
    AddNode(A, 1, 9.87);  
    AddNode(A, 2, 8.87);  
    AddNode(A, 3, 7.87);  
    AddNode(B, 4, 6.97);  
    AddNode(B, 5, 3.33);  
    AddNode(B, 6, 5.78);  
    PrintLL(A);  
    PrintLL(B);  
    Concatenate(A,B);  
    PrintLL(A);  
}
```

Results

Element 1: Balance of acc 3 is 7.87
Element 2: Balance of acc 2 is 8.87
Element 3: Balance of acc 1 is 9.87
End of the list.

Element 1: Balance of acc 6 is 5.78
Element 2: Balance of acc 5 is 3.33
Element 3: Balance of acc 4 is 6.97
End of the list.

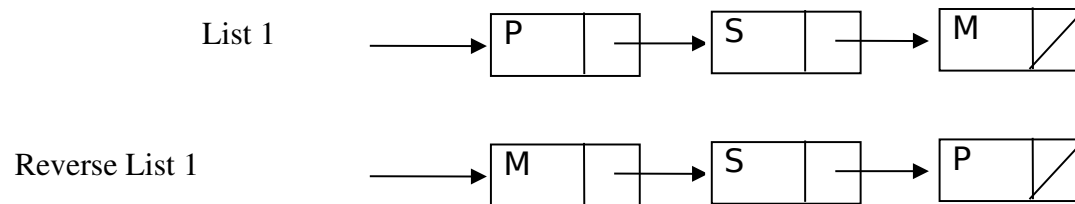
join

Element 1: Balance of acc 3 is 7.87
Element 2: Balance of acc 2 is 8.87
Element 3: Balance of acc 1 is 9.87
Element 4: Balance of acc 6 is 5.78
Element 5: Balance of acc 5 is 3.33
Element 6: Balance of acc 4 is 6.97
End of the list.

Reverse

Invert the order of all the elements.

The pointer to the last element becomes the list,
the the pointer to the list becomes the last
element...



Reverse

Many ways of achieving that...

E.g., two methods

Method 1: create a new LL, scan once to find how many elements, copy the last one, and keep adding nodes and scanning again. Copy address and delete the original.

Method 2: Scan once, swap the contents of the last element with the first one, keep going until swap the middle elements.



Reverse

We need code to search by position:

```
Node * SearchByPosition(Node *listpointer, int x) {  
    Node *current;  
    current = listpointer;  
    int pos=0;  
    while (true) {  
        if (current == NULL) { break; }  
        if (pos == x) { return current; }  
        current = current->next;  
        pos++;  
    }  
    printf("There are only %d elements in this list\n", x); return NULL;  
}
```


Reverse method 1

```
void ReverseLL1(Node * &listpointer) {
    Node *current;
    Node *prev;
    Node *reversedcopy=NULL;
    current = listpointer;
    int numbelements=0;
    while (true) {//scan once
        if (current == NULL) { break; }
        prev = current;
        current = current->next;
        numbelements++;
    }
    if(numbelements!=0){
        for(int count=0;count<numbelements;count++){
            Node *temp=SearchByPosition(listpointer, count);//find contents
            if(temp!=NULL) AddNode(reversedcopy,temp->accnumber,temp->balance);//copy contents
        }
        listpointer=reversedcopy;
        return;
    }
    else {
        printf("the list is empty, nothing to reverse\n");
        return;
    }
}
```

Reverse method 1

```
int main() {  
    A = NULL;..  
    AddNode(A, 1, 9.87);  
    AddNode(A, 2, 8.87);  
    AddNode(A, 3, 7.87);  
    AddNode(A, 4, 6.97);  
    PrintLL(A);  
    ReverseLL1(A);  
    PrintLL(A);  
}
```

Element 1: Balance of acc 4 is 6.97
Element 2: Balance of acc 3 is 7.87
Element 3: Balance of acc 2 is 8.87
Element 4: Balance of acc 1 is 9.87
End of the list.

the list is reversed

Element 1: Balance of acc 1 is 9.87
Element 2: Balance of acc 2 is 8.87
Element 3: Balance of acc 3 is 7.87
Element 4: Balance of acc 4 is 6.97
End of the list

Reverse method 1

Question 1: what is missing in ReverseLL1?

This is known as a ***memory leaking*** problem...

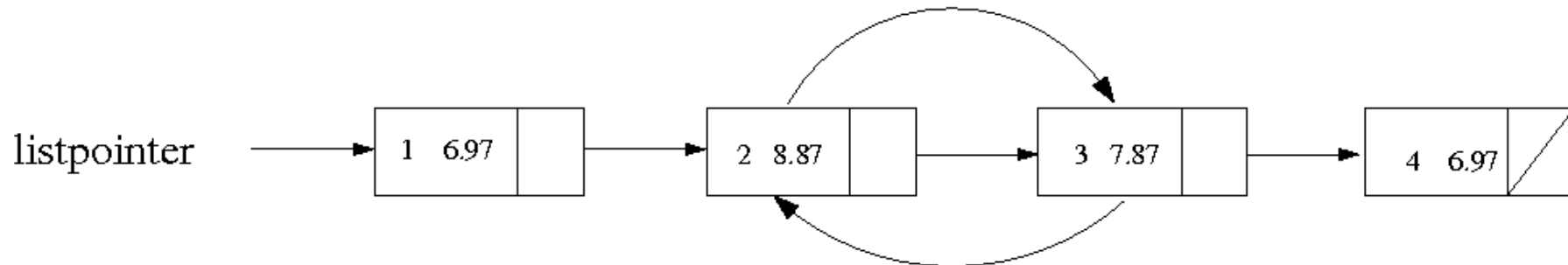
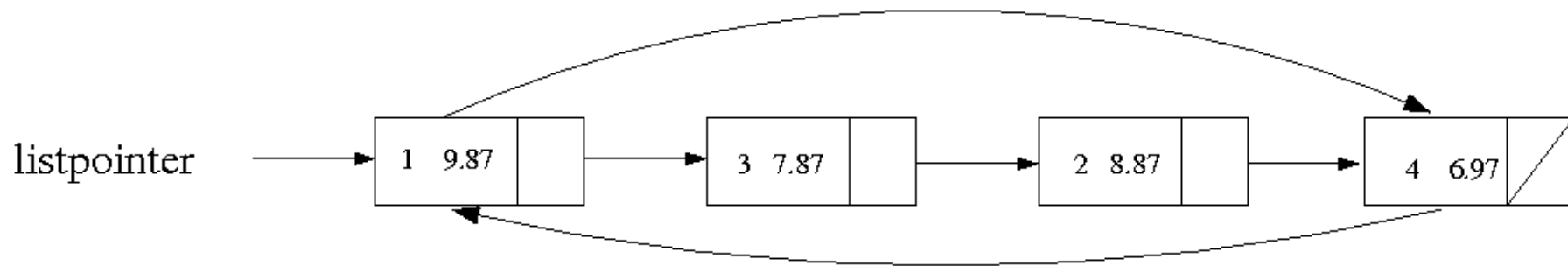
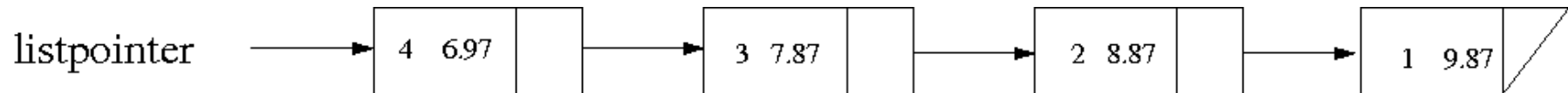
Question 2: is that efficient? How many scans/access do we have to do?



Reverse method 2

```
void ReverseLL2(Node * listpointer) {//Note we don't need & here
    Node *current, *prev, *temp, *temp2;
    current = listpointer;
    int numbelements=0;
    while (true) {//scan once
        if (current == NULL) { break; }
        prev = current;
        current = current->next;
        numbelements++;
    }
    if(numbelements!=0){
        for(int count=0;count<numbelements/2;count++){
            temp=SearchByPosition(listpointer, count);
            temp2=SearchByPosition(listpointer, numbelements-1-count);
            //swap values
            int accnumber_temp=temp->accnumber;
            float balance_temp=temp->balance;
            temp->accnumber = temp2->accnumber;
            temp->balance = temp2->balance;
            temp2->accnumber = accnumber_temp;
            temp2->balance = balance_temp;
        }
        printf("the list is reversed\n");
        return;
    }
    else { printf("the list is empty, nothing to reverse\n"); return; }
}
```

Reverse method 2



Circular Linked-lists

The last element points back to the first one

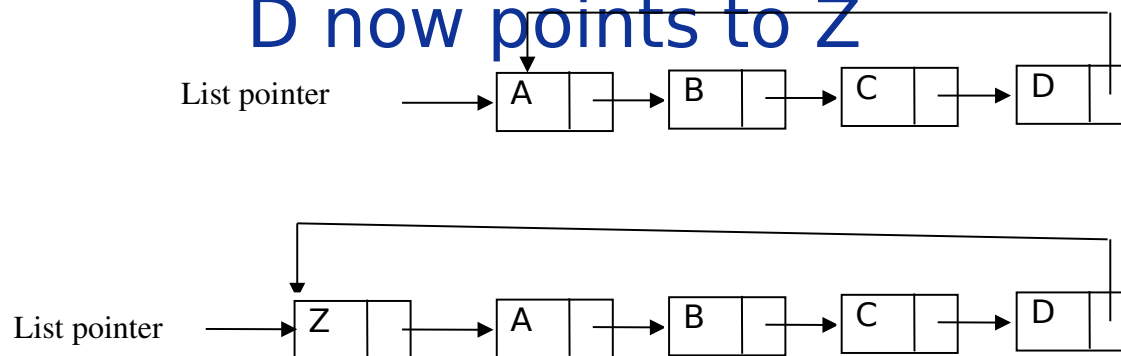
Adding nodes to the **middle** is easy...

Adding nodes to the **beginning** or **end** needs a different operation.

E.g., if we add Z to the start of the LL as we did before:

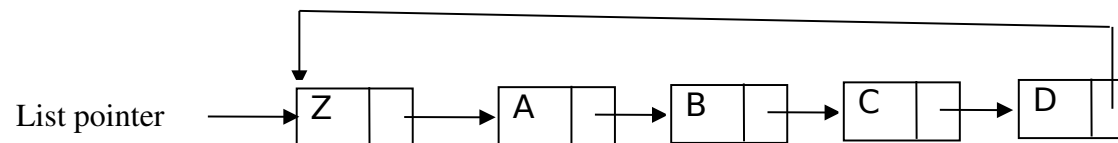
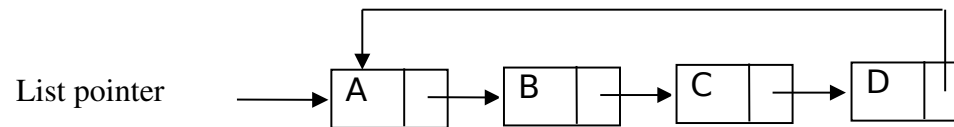
Z points to A

D now points to Z



Circular Linked-lists

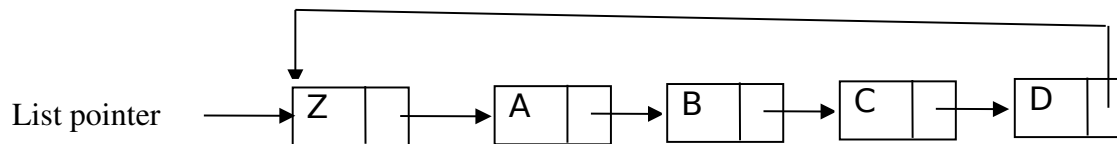
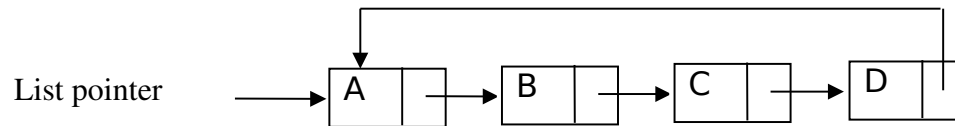
Question: When printing the linked-list, how do you know you reached the end?



Circular Linked-lists

Question: When printing the linked-list, how do you know you reached the end?

Answer: `current->next == listpointer`



Circular Linked-lists

Exercise:

Modify the code for the function

```
void PrintLL(Node *listpointer)
```

So it prints a circular Linked-list without looping forever



Circular Linked-lists

Solution:

```
void CLL_PrintLL(Node *listpointer) {  
    // print all elements  
    Node *current;  
    current = listpointer;  
    int element=1;  
    while (current->next != listpointer) {  
        printf("Element %d: Balance of acc %i is %1.2f\n",  
            element, current->accnumber, current->balance);  
        current = current->next;  
        element++;  
    }  
    //print last element  
    printf("Element %d: Balance of acc %i is %1.2f\n",  
        element, current->accnumber, current->balance);  
    printf("End of the list.\n");  
}
```

Doubly-linked Linked-lists

Have two pointers:

Forward (next)

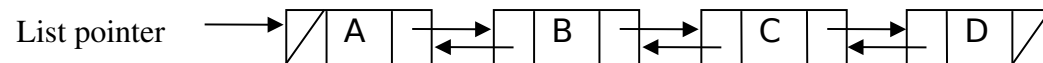
Backward (previous)

Advantages/disadvantages:

Search backwards, deal with neighbours simultaneously

More space in memory for the same amount of data

Operations have two pointers to update

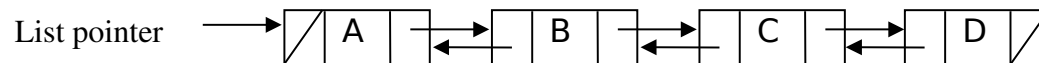


Doubly-linked Linked-lists

```
#include <stdio.h>

struct Node { //declaration
    int accnumber;
    float balance;
    Node *next;
    Node *previous;
};

Node *A, *B; //declaration
```



Challenge:

1) Can you think of a better method to reverse a linked-list, without any allocation or deallocation?

